

# Exact and efficient simulation of photon quantum interference using tensor train decomposition

**S. Fldzhyan**<sup>1,2</sup>, G. Ryzhakov<sup>3</sup>, M. Saygin<sup>2,4,5</sup>, I. Oseledets<sup>3,6</sup>, and S. Straupe<sup>2,4</sup>

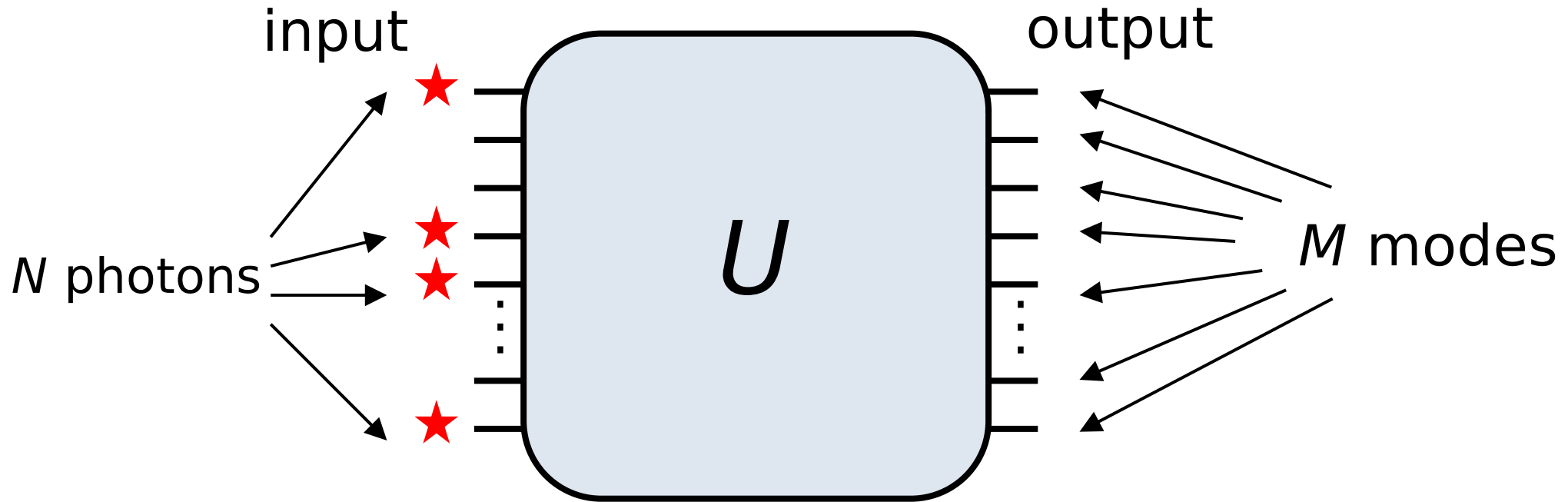
<sup>1</sup>Russian Quantum Center, <sup>2</sup>MSU Quantum Technology Center, <sup>3</sup>Skolkovo Institute of Science and Technology, <sup>4</sup>Sber Quantum Technology Center, <sup>5</sup>Laboratory of Quantum Engineering of Light South Ural State University, <sup>6</sup>Artificial Intelligence Research Institute (AIRI)

**Russian Supercomputing Days 2024**

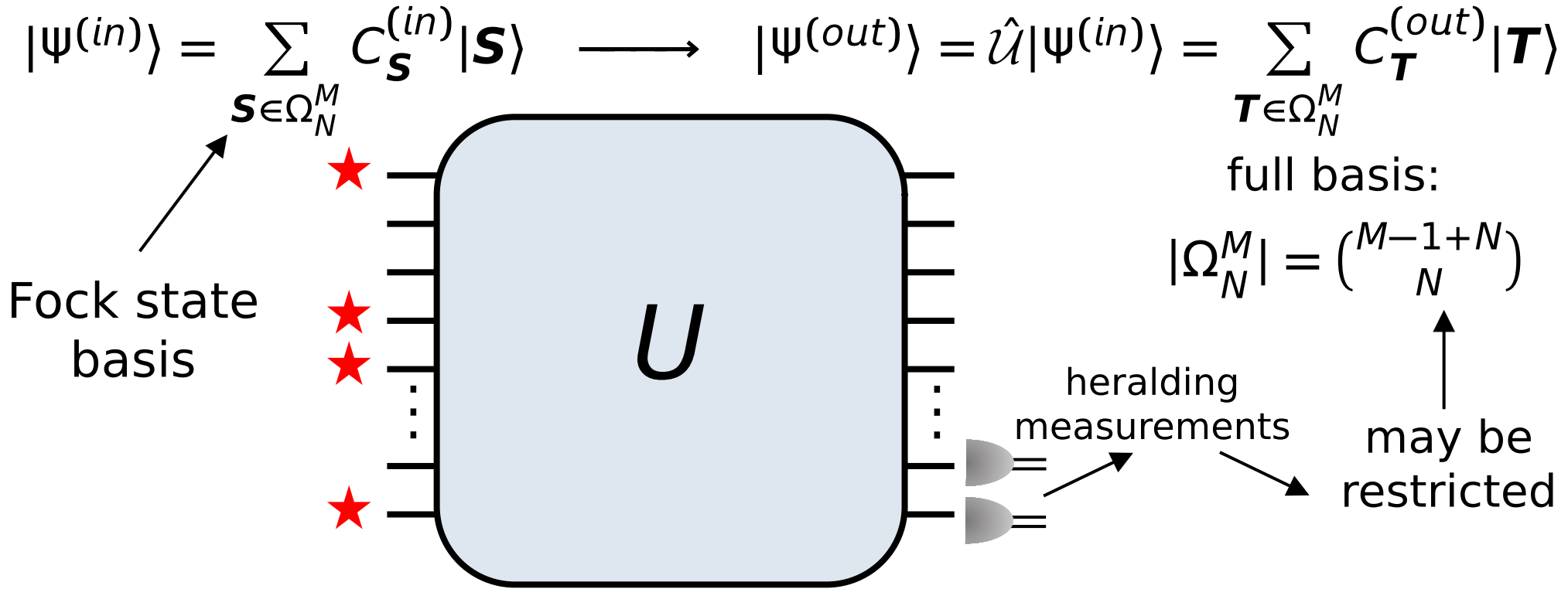
September 24

# Problem under consideration

A standart setup of linear optical quanutm computation



# Problem under consideration



The task: precisely compute  $C_{\mathbf{T}}^{(out)}$  for a lot of different  $U$

# Fock state

occupation numbers

assignment list:  $1 \leq s_1 \leq s_2 \leq \dots \leq s_N \leq M$

$$|\mathbf{S}\rangle = |S_1, S_2, S_3, \dots, S_M\rangle = |\mathbf{s}\rangle = |\{s_1, s_2, s_3, \dots, s_N\}\rangle$$

$S_i$  photons  
in mode  $i$

$$\sum_{i=1}^M S_i = N$$

" $i$ -th photon  
in mode  $s_i$ "

$$\sum_{j=1}^N \delta_{is_j} = S_i$$

examples:

full basis:

$$|\Omega_N^M| = \binom{M-1+N}{N}$$

$$|1, 2, 0, 1\rangle = |\{1, 2, 2, 4\}\rangle$$

$$|0, 2, 4\rangle = |\{2, 2, 3, 3, 3, 3\}\rangle$$

# Fock operators

$\Omega_N^M \xrightarrow{\hat{a}_i^\dagger} \Omega_{N+1}^M$  ← creation operator      annihilation operator →  $\Omega_N^M \xrightarrow{\hat{a}_i} \Omega_{N-1}^M$

$[\hat{a}_i, \hat{a}_j^\dagger] = \delta_{ij}$

Definition:  $[\hat{a}_i, \hat{a}_j] = 0, [\hat{a}_i^\dagger, \hat{a}_j^\dagger] = 0$

$\hat{a}_i^\dagger |\dots, S_i, \dots\rangle = \sqrt{S_i + 1} |\dots, S_i + 1, \dots\rangle$        $\hat{a}_i |\dots, S_i, \dots\rangle = \sqrt{S_i} |\dots, S_i - 1, \dots\rangle$

$\hat{a}^\dagger = \begin{pmatrix} 0 & 0 & \dots & & & \\ 1 & 0 & \dots & & & \\ 0 & \sqrt{2} & \dots & & & \\ \vdots & \vdots & \ddots & & & \\ \vdots & \vdots & \ddots & \vdots & & \\ \dots & \dots & \sqrt{K} & 0 & & \\ \dots & \dots & 0 & \sqrt{K+1} & & \end{pmatrix} \begin{matrix} \Omega_K^1 \\ \updownarrow \Omega_{K+1}^1 \end{matrix}$

$\hat{a} = \begin{pmatrix} 0 & 1 & 0 & \dots & & \\ 0 & 0 & \sqrt{2} & \dots & & \\ \vdots & \vdots & \vdots & \ddots & & \\ \dots & \dots & \dots & \sqrt{K-1} & 0 & \\ \dots & \dots & 0 & \sqrt{K} & & \end{pmatrix} \begin{matrix} \Omega_K^1 \\ \updownarrow \Omega_{K-1}^1 \end{matrix}$

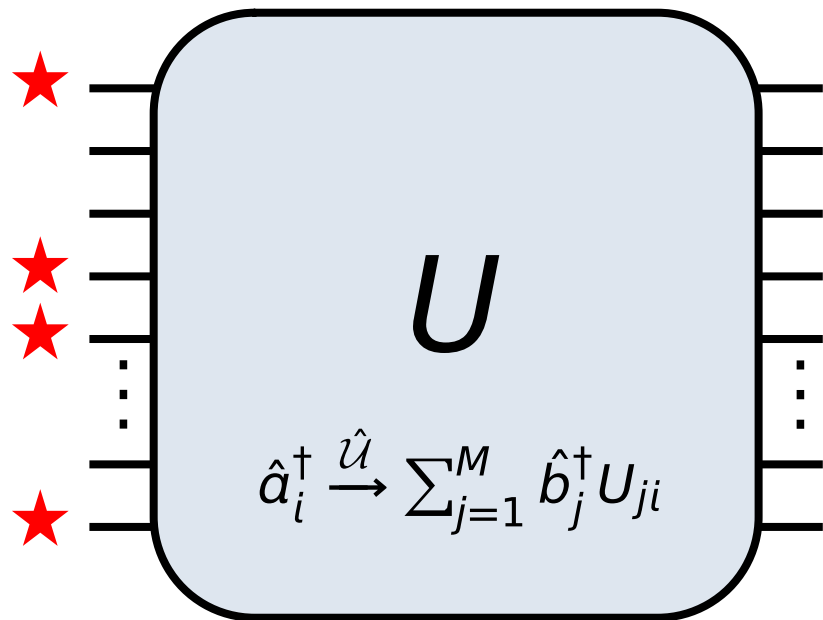
very sparse

# Fock state evolution

$$|\mathbf{s}\rangle = |S_1, S_2, S_3 \dots, S_M\rangle = \prod_{i=1}^M \frac{(\hat{a}_i^\dagger)^{S_i}}{\sqrt{S_i!}} |\text{vac}\rangle = |\mathbf{s}\rangle = |\{s_1, s_2, s_3 \dots, s_N\}\rangle = \frac{1}{\sqrt{\mathbf{s}!}} \prod_{i=1}^N \hat{a}_{s_i}^\dagger |\text{vac}\rangle$$

$\mathbf{s}! = \prod_{i=1}^M S_i!$

transformation:



output state

$$|\mathbf{s}\rangle = \frac{1}{\sqrt{\mathbf{s}!}} \prod_{i=1}^N \hat{a}_{s_i}^\dagger |\text{vac}\rangle \leftarrow \text{input state}$$

$$|\psi^{(out)}\rangle = \frac{1}{\sqrt{\mathbf{s}!}} \prod_{i=1}^N \left( \sum_{j=1}^M \hat{b}_j^\dagger U_{js_i} \right) |\text{vac}\rangle = \sum_{\mathbf{t} \in \Omega_N^M} U_{\mathbf{t}\mathbf{s}} |\mathbf{t}\rangle$$

transformation amplitude  $\rightarrow U_{\mathbf{t}\mathbf{s}} = \frac{\text{perm}[U^{\mathbf{t}\mathbf{s}}]}{\sqrt{\mathbf{s}!\mathbf{t}!}} \quad O(N2^N)$

$$[U^{\mathbf{t}\mathbf{s}}]_{ij} \equiv U_{t_i s_j}, \text{perm}[A] \equiv \sum_{\sigma \in \mathcal{S}_N} \prod_{i=1}^N A_{\sigma_i i}$$

Total complexity of a problem is  $\sim O\left(N2^N \binom{M-1+N}{N}\right)$

# SLOS method

$$\frac{1}{\sqrt{\mathbf{s}!}} \prod_{i=1}^N \left( \sum_{j=1}^M \hat{b}_j^\dagger U_{js_i} \right) |vac\rangle = \frac{1}{\sqrt{\mathbf{s}!}} \left( \sum_{j=1}^M \hat{b}_j^\dagger U_{js_1} \right) \cdots \left( \sum_{j=1}^M \hat{b}_j^\dagger U_{js_{N-1}} \right) \cdot \left( \sum_{j=1}^M \hat{b}_j^\dagger U_{js_N} \right) |vac\rangle = |\psi^{(out)}\rangle$$

Mat · Vec : nnz [Mat] multiplications

Matrices have a lot of zeroes — total complexity of a problem is  $O\left(N^{\binom{M-1+N}{N}}\right)$

Generalization to arbitrary set of input and output Fock bases is unclear

# TT decomposition

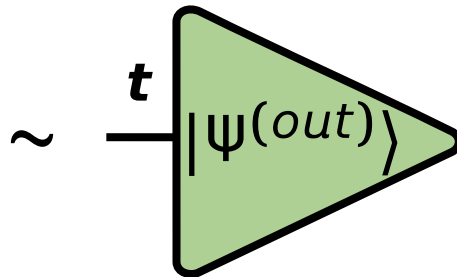
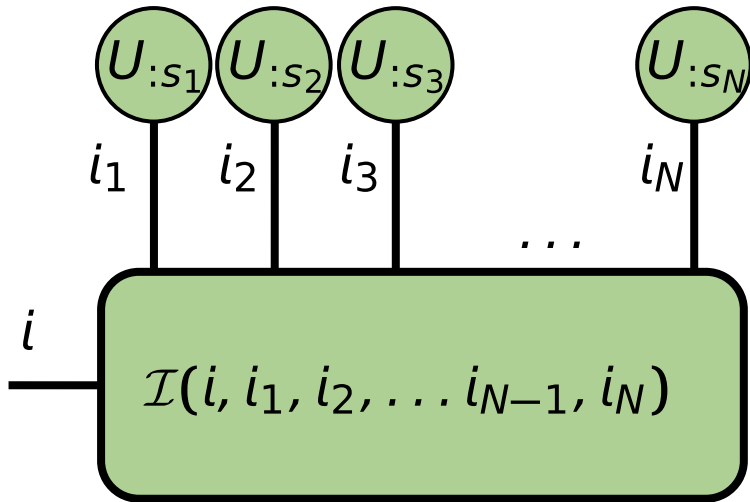
may not be a full basis



Suppose output basis is given by a set of assignment lists:  $L[i] = \{t_1, t_2, \dots, t_N\}$

A key observation:

$$\text{perm}[U^{\mathbf{t}\mathbf{s}}] = \sum_{\sigma \in S_N} \prod_{j=1}^N U_{t_{\sigma_j} s_j} \stackrel{\mathbf{t} = L[i]}{=} \mathbf{t}! \sum_{i_1=1}^M \sum_{i_2=1}^M \cdots \sum_{i_N=1}^M \mathcal{I}(i, i_1, i_2, \dots, i_N) \prod_{k=1}^N U_{i_k s_k}$$



Indicator tensor:

$$\mathcal{I}(i, i_1, \dots, i_N) = \begin{cases} 1, & \{i_1, \dots, i_N\} = L[i] \\ 0, & \text{otherwise} \end{cases}$$



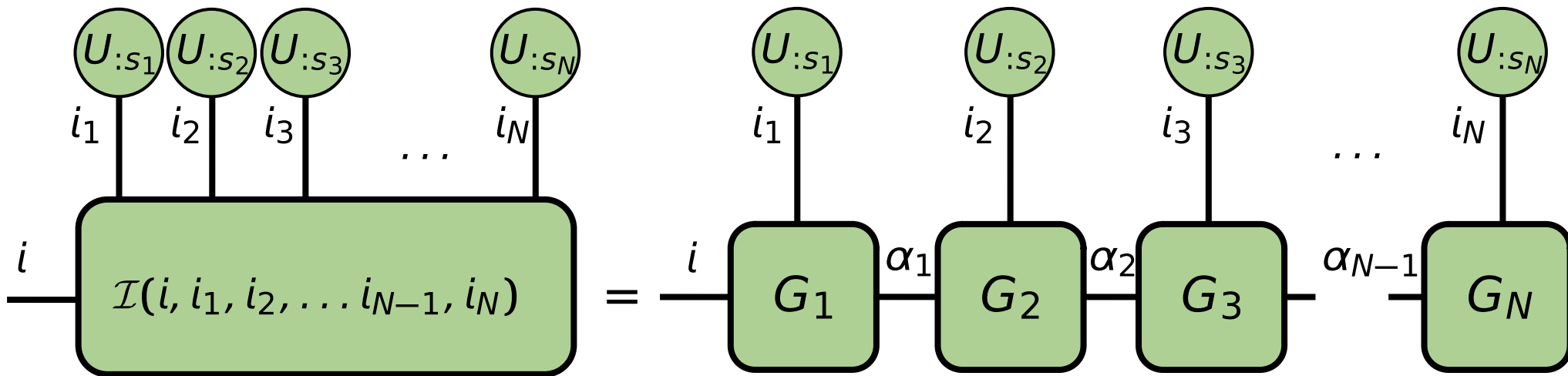
$\mathcal{I}$  is completely defined by  $L$



# TT decomposition

$$\sum_{i_1=1}^M \sum_{i_2=1}^M \cdots \sum_{i_N=1}^M \mathcal{I}(i, i_1, i_2, \dots, i_N) \prod_{k=1}^N U_{i_k s_k} = \sum_{\alpha_1 \alpha_2 \dots \alpha_{N-1}} \prod_{k=1}^N \left( \sum_{i_k=1}^M G_k(\alpha_{k-1}; i_k; \alpha_k) U_{i_k s_k} \right)$$

cores of TT



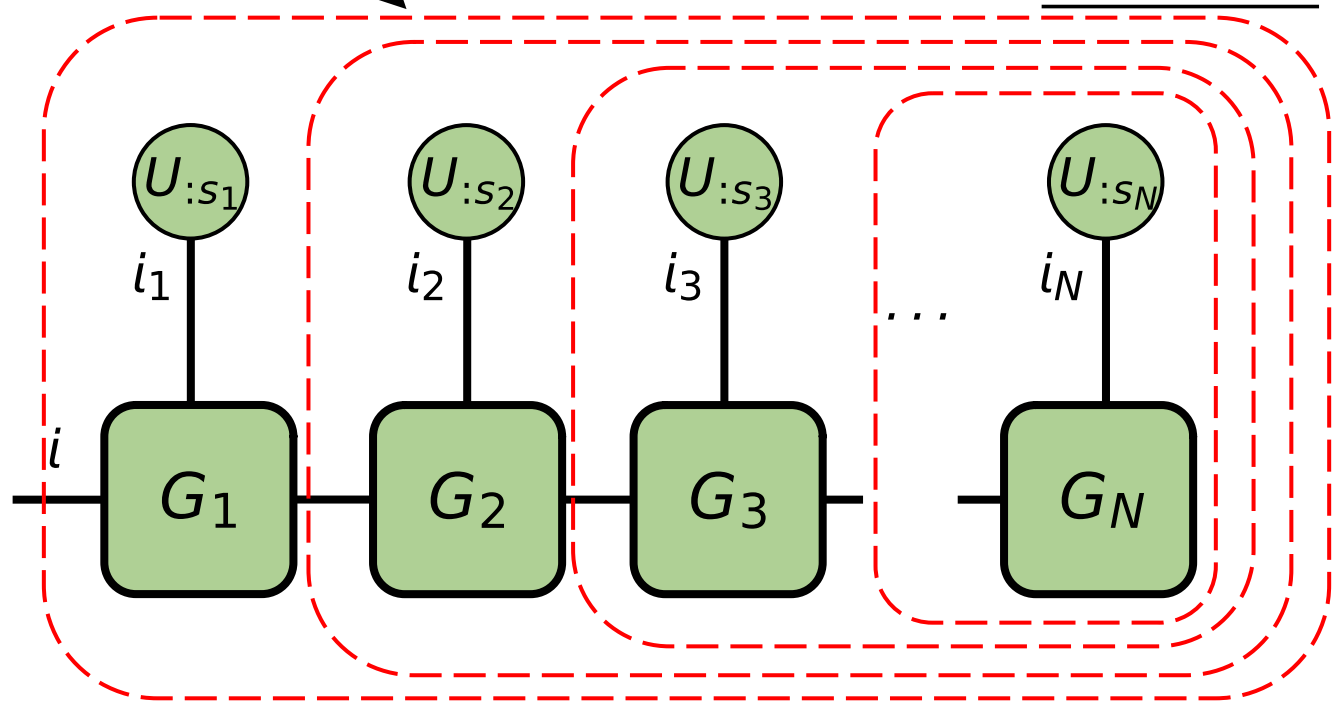
Derivative function decomposition method introduced in

Ryzhakov, G., & Oseledets, I. (2023). «Constructive TT-representation of the tensors given as index interaction functions with applications». The Eleventh International Conference on Learning Representations

# Single input, arbitrary output

A universal recipe for single Fock input and arbitrary Fock basis output

order of evaluation  
↙



only matrix by vector multiplication

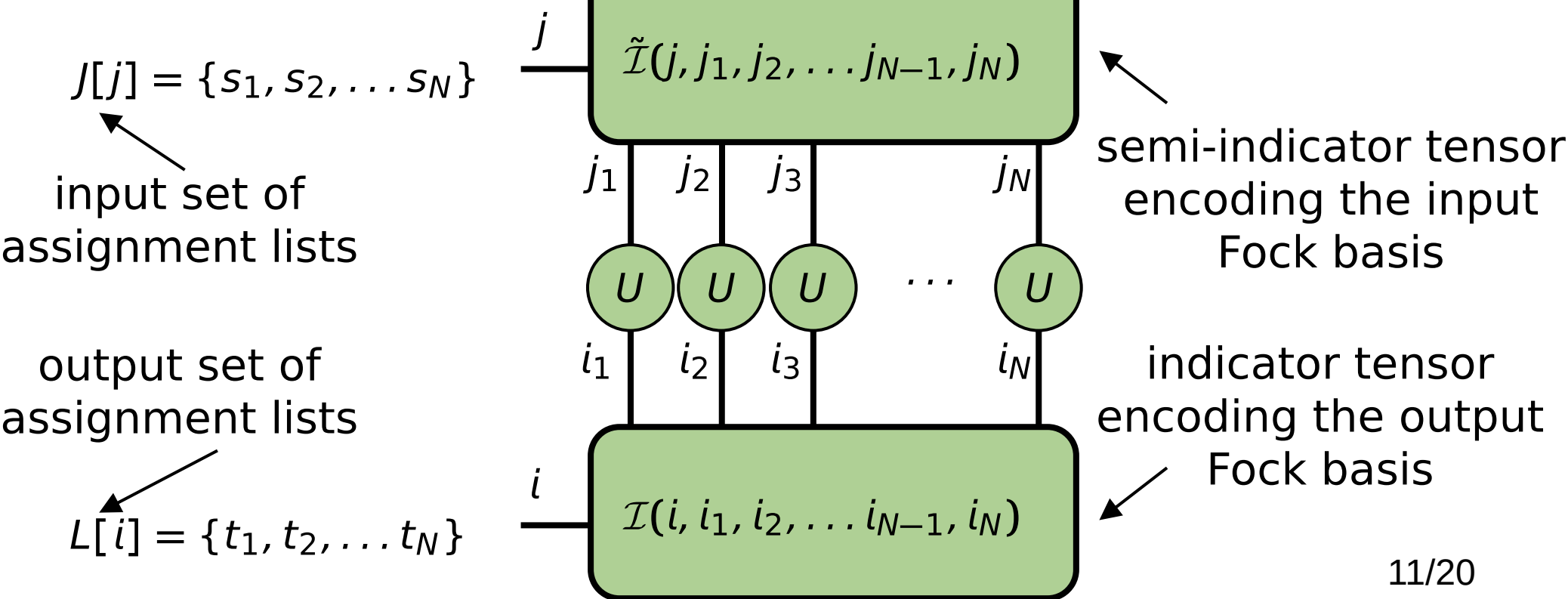
input:  $\mathbf{s}$       output:  $L[i]$

for full output basis:  
 $O\left(N^{\binom{M-1+N}{N}}\right)$

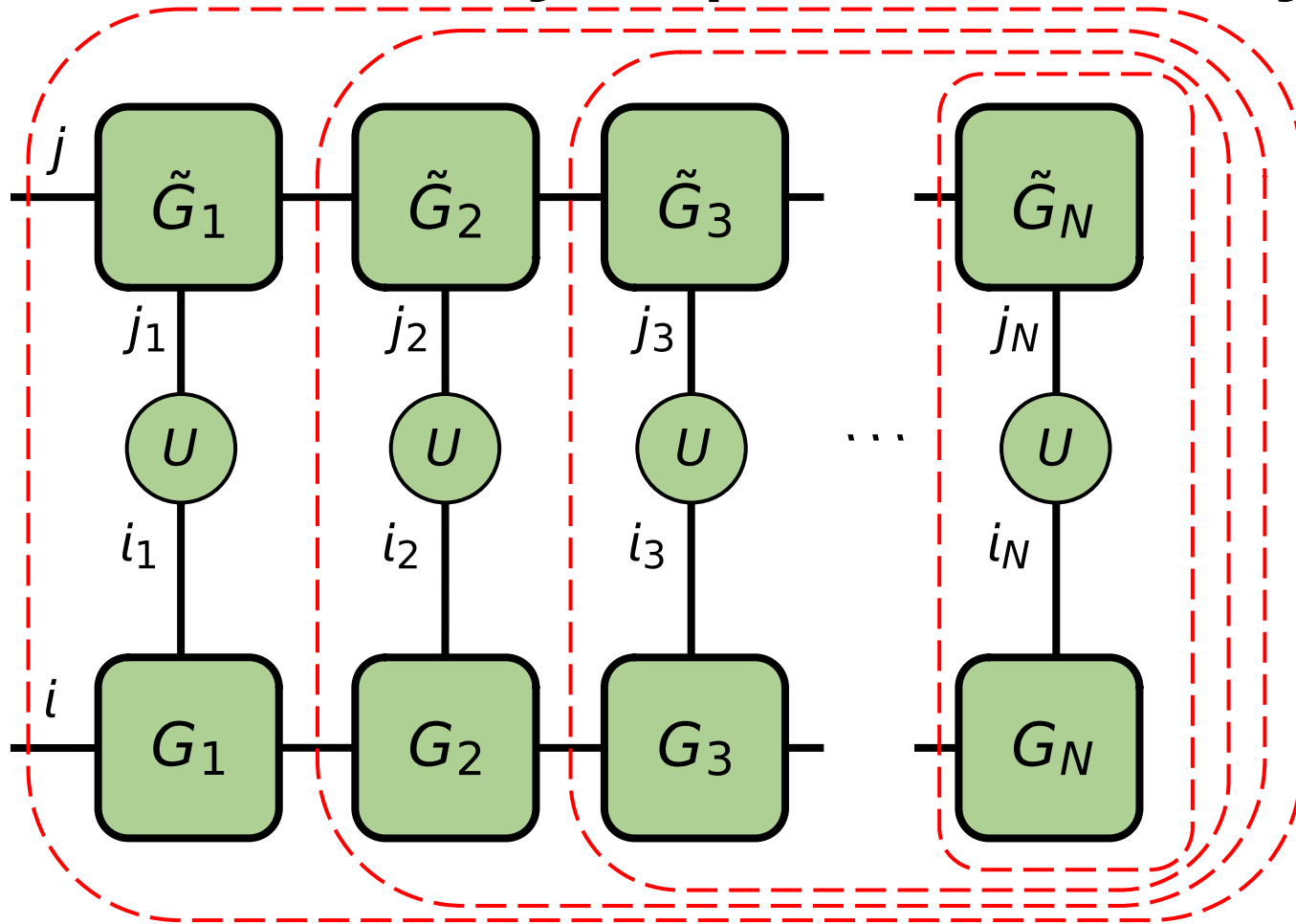
↓  
equivalent to SLOS

# Arbitrary input, arbitrary output

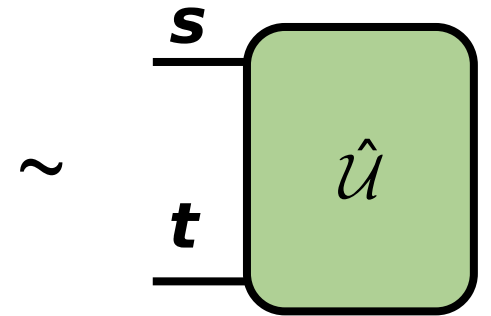
$$U_{\mathbf{ts}} \sim \sum_{i_1 j_1=1}^M \sum_{i_2 j_2=1}^M \cdots \sum_{i_N j_N=1}^M \mathcal{I}(i, i_1, i_2, \dots, i_N) \tilde{\mathcal{I}}(j, j_1, j_2, \dots, j_N) \prod_{k=1}^N U_{i_k j_k}$$



# Arbitrary input, arbitrary output



matrix of transformation  
amplitudes



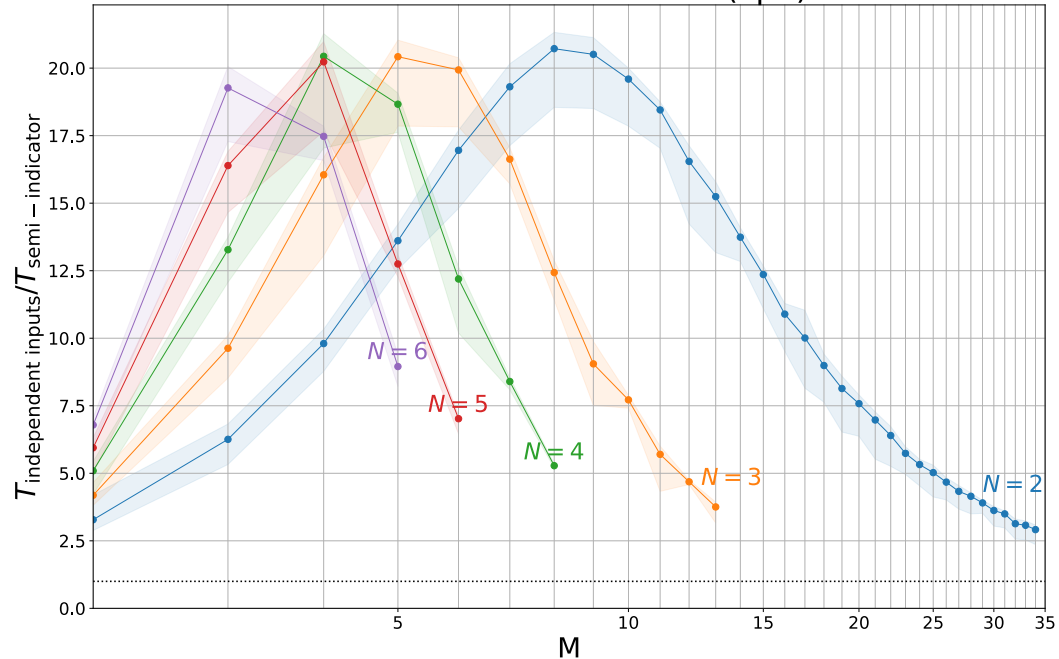
for full output  
and input basis:

$$O\left(N \binom{M-1+N}{N}^2\right)$$

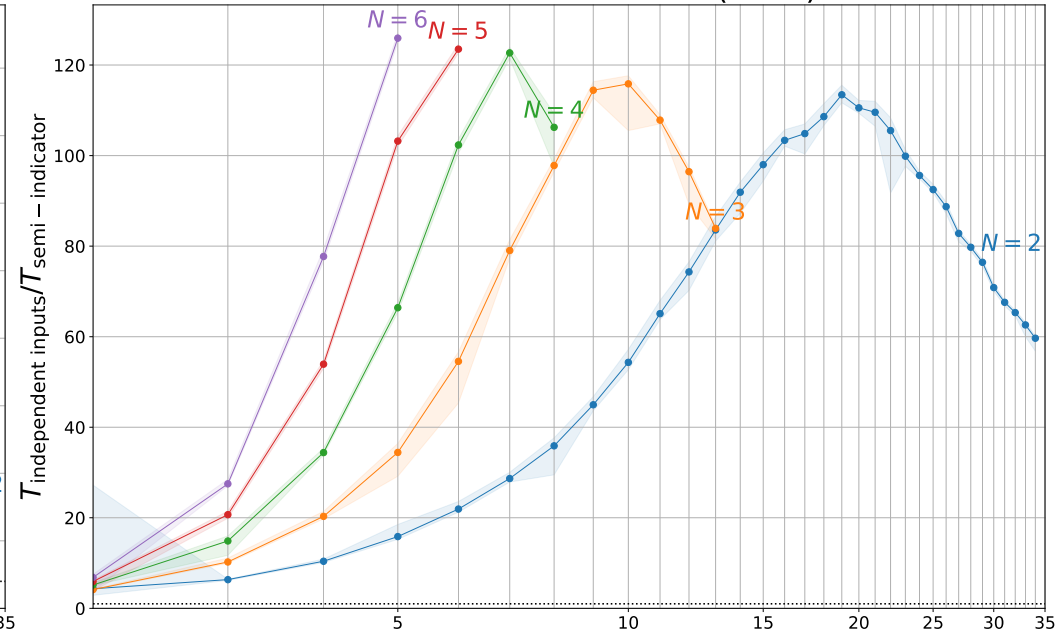
# Numerical comparison

How faster is it to calculate transformation amplitudes with the semi-indicator tensor than to calculate each input separately for full bases?

calculation time relation (cpu)



calculation time relation (cuda)

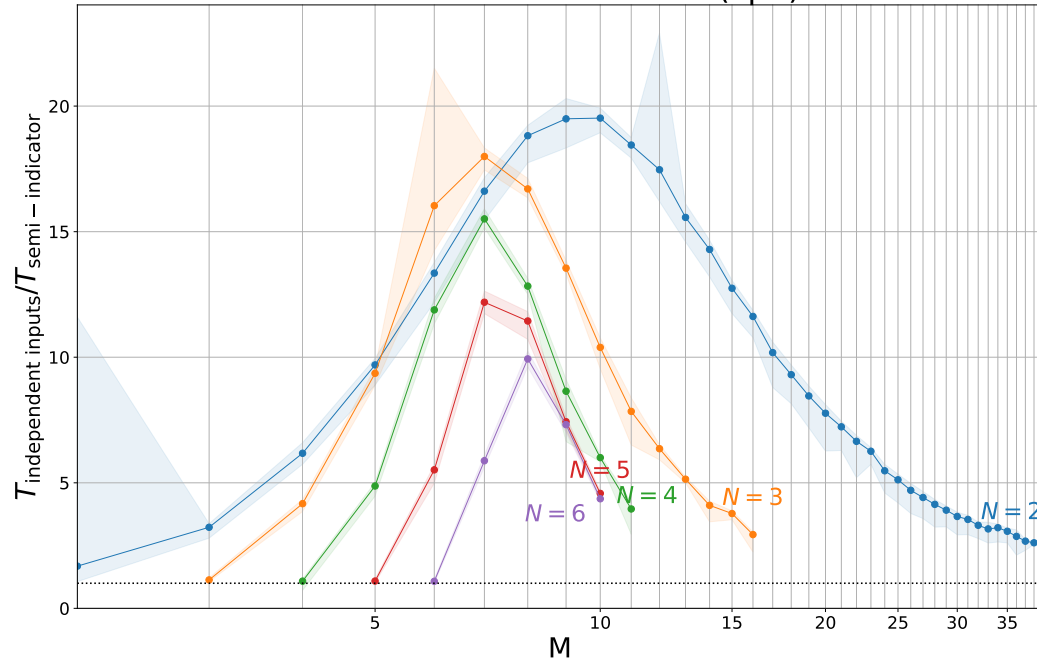


Both methods are asymptotically equivalent. For full bases:  $O\left(N \binom{M-1+N}{N}^2\right) \cdot 13/20$

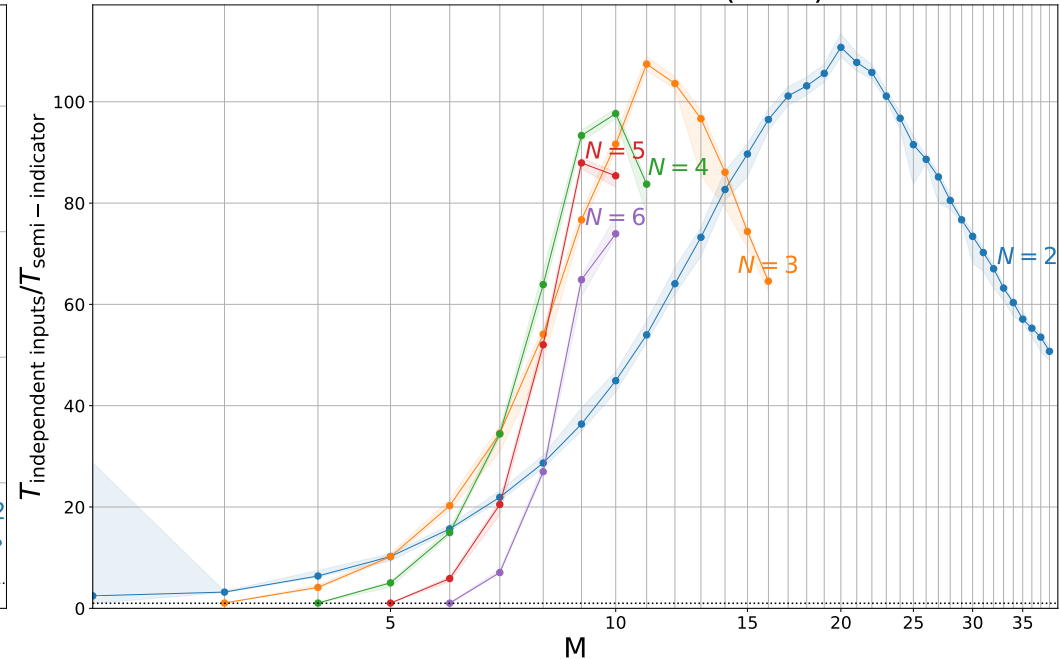
# Numerical comparison

How faster is it to calculate transformation amplitudes with the semi-indicator tensor than to calculate each input separately for unbunched bases (assignment lists have no repetitions)?

calculation time relation (cpu)

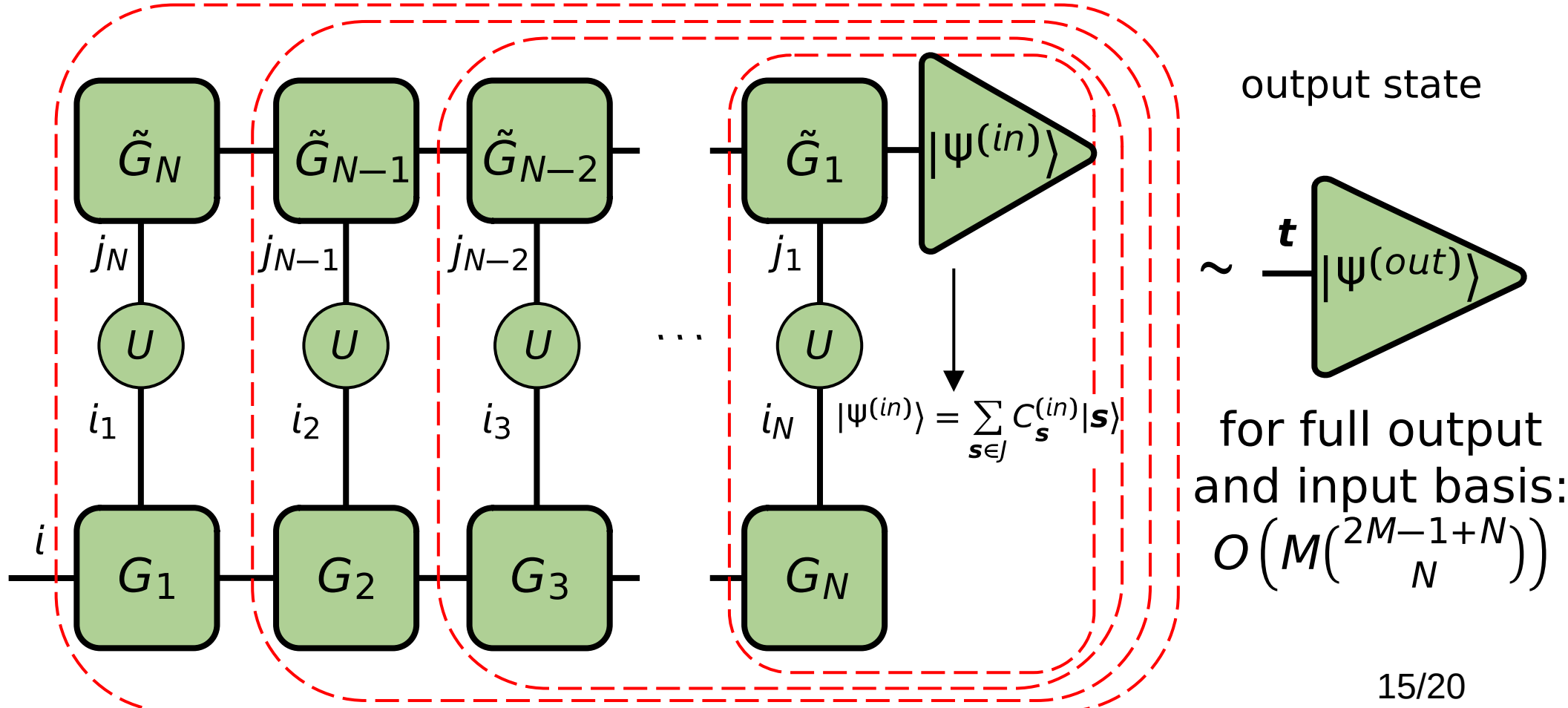


calculation time relation (cuda)



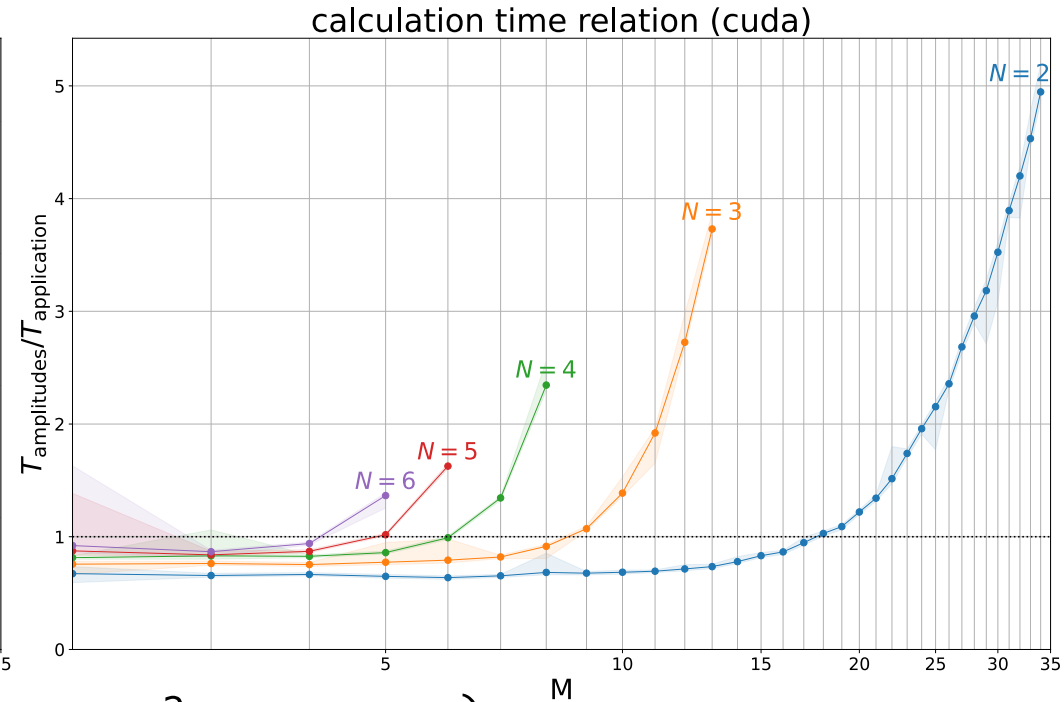
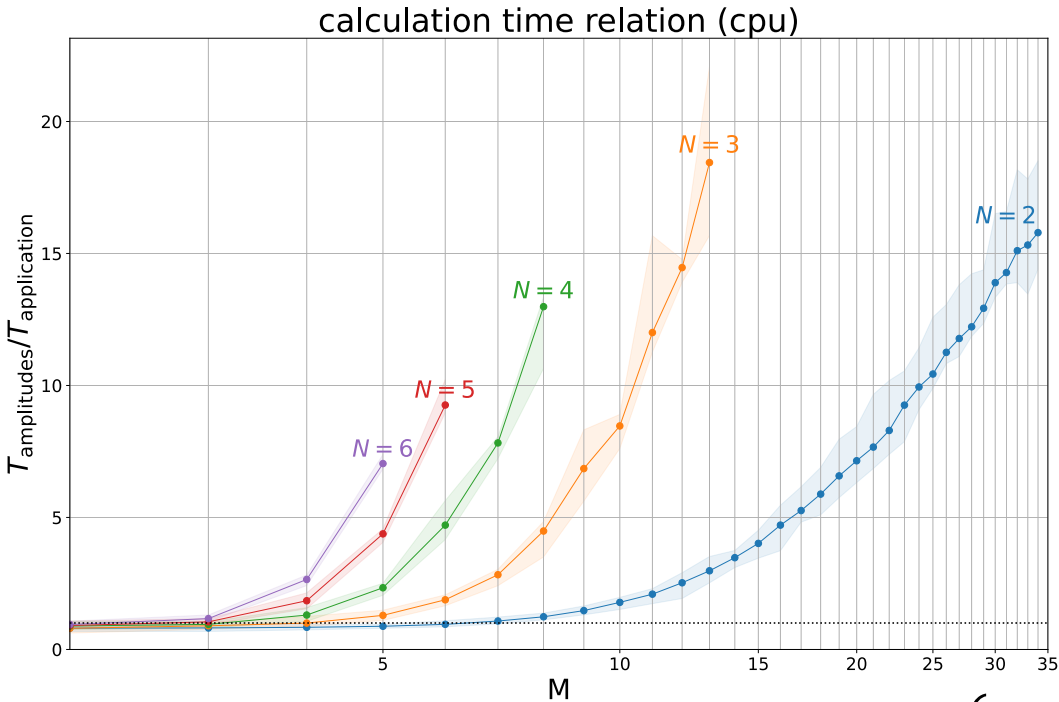
Both methods are asymptotically equivalent.

# Arbitrary input, arbitrary output



# Numerical comparison

How faster is it to calculate all the transformation amplitudes compared to calculating the application of the transformation for full bases?



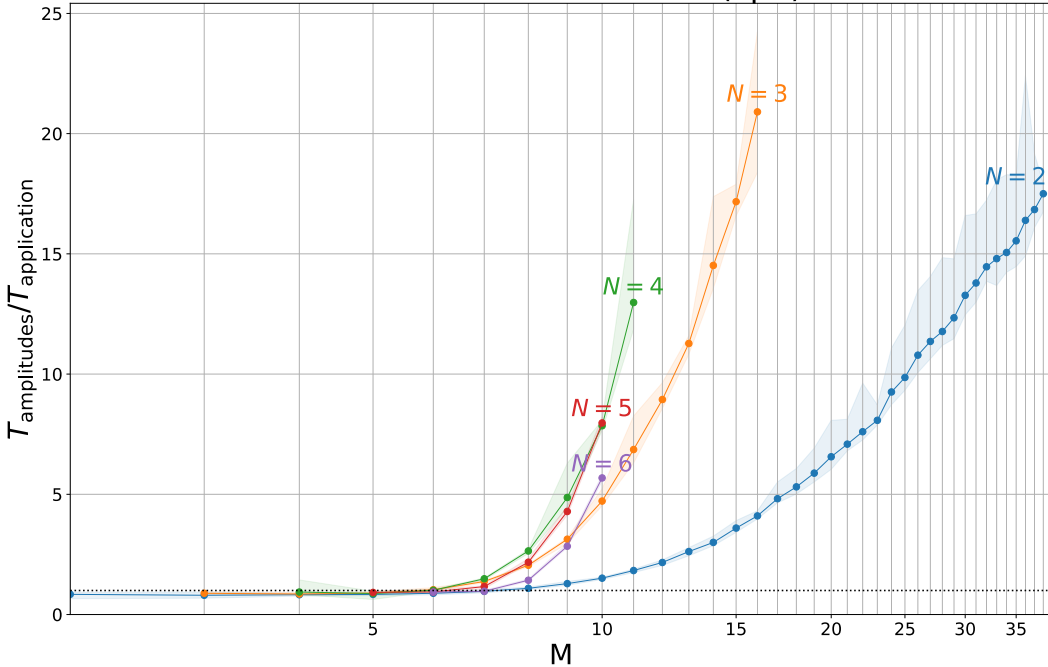
For full bases:  $O\left(N \binom{M-1+N}{N}^2 / M \binom{2M-1+N}{N}\right)$



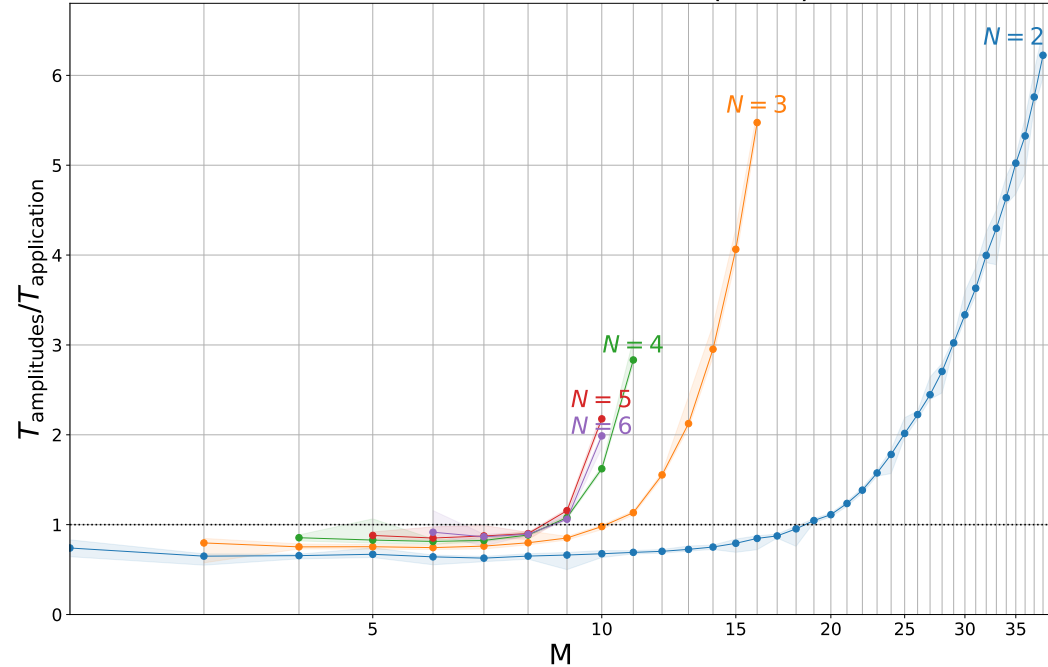
# Numerical comparison

How faster is it to calculate all the transformation amplitudes compared to calculating the application of the transformation for unbunched bases?

calculation time relation (cpu)



calculation time relation (cuda)



# Possible generalization

Other quantities involving a sum over permutations

Distinguishable photons:

$$S_{ij} = \langle \phi_i | \phi_j \rangle$$

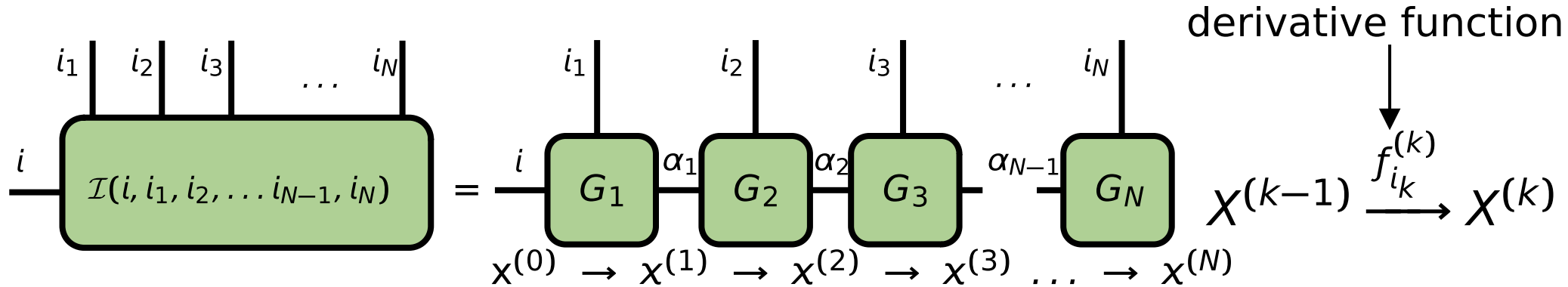
$$\mathcal{P}(\mathbf{t}|\mathbf{s}) = \sum_{\sigma, \tau \in \mathcal{S}_N} \prod_{i=1}^N S_{\sigma_i \tau_i} [\bar{U}^{\mathbf{t}\mathbf{s}}]_{i\sigma_i} [U^{\mathbf{t}\mathbf{s}}]_{i\tau_i}$$

Gaussian boson sampling:

$$\mathcal{P}(\mathbf{s}) = \frac{\text{haf}[A^{\mathbf{s}}]}{\sqrt{\det[\sigma]} \mathbf{s}!}$$

$$\text{haf}[A] = \frac{1}{M! 2^M} \sum_{\sigma \in \mathcal{S}_{2M}} \prod_{i=1}^M A_{\sigma_{2i-1} \sigma_{2i}}$$

# Derivative function decomposition



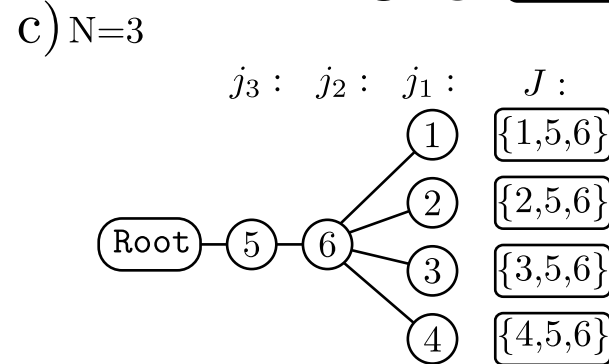
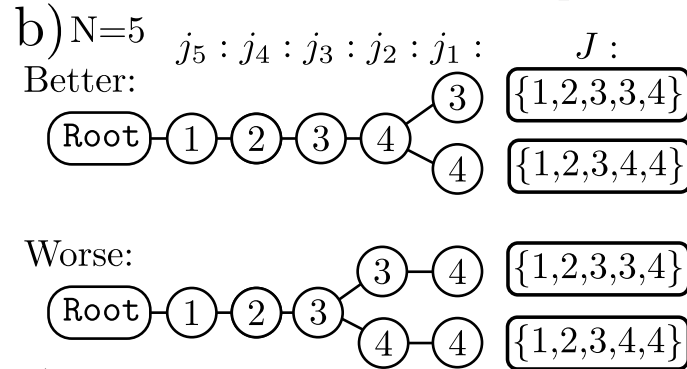
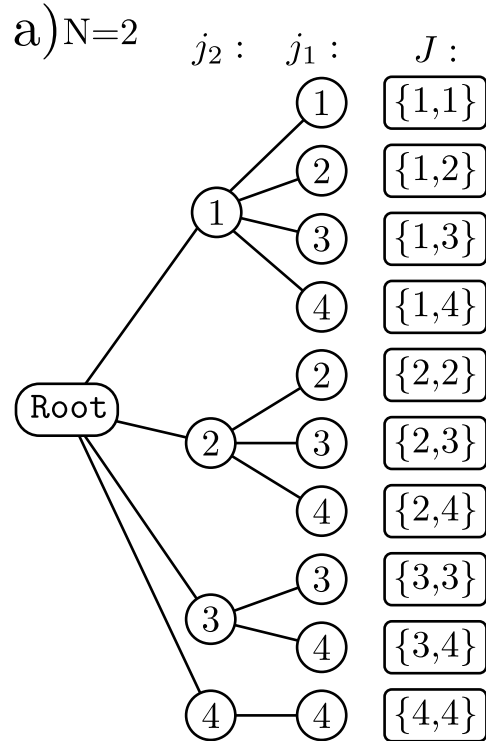
$$X^{(0)} \sim L, X^{(0)} = \{T_1, T_2, T_3, \dots, T_M\} \sim L[i]$$

$$\alpha_{k-1} = \overline{1, |X^{(k-1)}|}, \alpha_k = \overline{1, |X^{(k)}|}$$

for  $G_k(\alpha_{k-1}; i_k; \alpha_k) = 0$  or  $1; \forall 1 \leq k \leq N$ :

$$x^{(k)} = f_{i_k}^{(k)}(x^{(k-1)}) = \begin{cases} x^{(k-1)}[i_k] \leftarrow x^{(k-1)}[i_k] - 1; & x^{(k-1)}[i_k] > 0, \\ \text{None}, & \text{else} \end{cases}$$

# Derivative function decomposition



$$x^{(k)} = \tilde{f}_{j_k}^{(k)}(x^{(k-1)}) = \begin{cases} x^{(k-1)}.parent, & \text{if } x^{(k-1)}.key = j_k, \\ \text{None}, & \text{else} \end{cases}, \quad \forall 1 \leq k \leq N.$$